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U.S. DEPARTMENT OF COMMERCE PATENT AND TRADEMARK OFFICE

ATTORNEY'S DOCKET NUMBER

**TRANSMITTAL LETTER TO THE UNITED STATES
DESIGNATED/ELECTED OFFICE (DO/EO/US)
CONCERNING A FILING UNDER 35 U.S.C. 371**

032326-132

U.S. APPLICATION NO. (If known, see 37 C.F.R. 1.5)

Unassigned

09/807607

INTERNATIONAL APPLICATION NO.
PCT/FR99/02172INTERNATIONAL FILING DATE
13 September 1999PRIORITY DATE CLAIMED
16 October 1998

TITLE OF INVENTION

COUNTERMEASURE METHOD IN AN ELECTRONIC COMPONENT USING A SECRET KEY CRYPTOGRAPHIC ALGORITHM

APPLICANT(S) FOR DO/EO/US

Christophe CLAVIER and Olivier BENOIT

Applicant herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information:

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☒ This is an express request to begin national examination procedures (35 U.S.C. 371(f)) at any time rather than delay examination until the expiration of the applicable time limit set in 35 U.S.C. 371(b) and the PCT Articles 22 and 39(1).
4. ☒ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☒ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
 - a. ☒ is transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☒ has been transmitted by the International Bureau.
 - c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US).
6. ☒ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☒ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. 371(c)(3))
 - a. ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ have been transmitted by the International Bureau.
 - c. ☐ have not been made; however, the time limit for making such amendments has NOT expired.
 - d. ☒ have not been made and will not be made.
8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371(c)(3)).
9. ☐ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☐ A translation of the annexes to the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11. to 16. below concern other document(s) or information included:

11. ☒ An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☒ A FIRST preliminary amendment.
 ☐ A SECOND or SUBSEQUENT preliminary amendment.
14. ☐ A substitute specification.
15. ☐ A change of power of attorney and/or address letter.
16. ☐ Other items or information:

16 APR 2001

U.S. APPLICATION NO. (If known / see 37 CFR 1.50)

Unassigned

09/807607

INTERNATIONAL APPLICATION NO

PCT/FR99/02172

ATTORNEY'S DOCKET NUMBER

032326-132

17. ☒ The following fees are submitted:

CALCULATIONS

PTO USE ONLY

Basic National Fee (37 CFR 1.492(a)(1)-(5)):

Neither international preliminary examination fee (37 CFR 1.482)
nor international search fee (37 CFR 1.445(a)(2)) paid to USPTO
and International Search Report not prepared by the EPO or JPO \$1,000.00 (960)

International preliminary examination fee (37 CFR 1.482) not paid to
USPTO but International Search Report prepared by the EPO or JPO \$860.00 (970)

International preliminary examination fee (37 CFR 1.482) not paid to USPTO
but international search fee (37 CFR 1.445(a)(2)) paid to USPTO \$710.00 (958)

International preliminary examination fee paid to USPTO (37 CFR 1.482)
but all claims did not satisfy provisions of PCT Article 33(1)-(4) \$690.00 (956)

International preliminary examination fee paid to USPTO (37 CFR 1.482)
and all claims satisfied provisions of PCT Article 33(1)-(4) \$100.00 (962)

ENTER APPROPRIATE BASIC FEE AMOUNT =

\$ 860.00

Surcharge of \$130.00 (154) for furnishing the oath or declaration later than
months from the earliest claimed priority date (37 CFR 1.492(e)).

20 ☐ 30 ☐

\$ -0-

Claims

Number Filed

Number Extra

Rate

Total Claims

16 -20 =

-0-

X\$18.00 (966)

\$ -0-

Independent Claims

3 -3 =

-0-

X\$80.00 (964)

\$ -0-

Multiple dependent claim(s) (if applicable)

+ \$270.00 (968)

\$ -0-

TOTAL OF ABOVE CALCULATIONS =

\$ 860.00

Reduction for 1/2 for filing by small entity, if applicable (see below).

\$ -0-

SUBTOTAL =

\$ 860.00

Processing fee of \$130.00 (156) for furnishing the English translation later than
months from the earliest claimed priority date (37 CFR 1.492(f)).

20 ☐ 30 ☐

\$ -0-

+

TOTAL NATIONAL FEE =

\$ -0-

Fee for recording the enclosed assignment (37 CFR 1.21(h)). The assignment must be accompanied by
an appropriate cover sheet (37 CFR 3.28, 3.31). \$40.00 (581) per property +

\$ -0-

TOTAL FEES ENCLOSED =

\$ 860.00

Amount to be:
refunded

\$

charged

\$

a. ☐ Small entity status is hereby claimed.b. ☒ A check in the amount of \$ 860.00 to cover the above fees is enclosed.c. ☐ Please charge my Deposit Account No. 02-4800 in the amount of \$ to cover the above fees. A duplicate copy of this sheet is enclosed.d. ☐ The Commissioner is hereby authorized to charge any additional fees which may be required, or credit any overpayment to Deposit Account No. 02-4800. A duplicate copy of this sheet is enclosed.

NOTE: Where an appropriate time limit under 37 CFR 1.494 or 1.495 has not been met, a petition to revive (37 CFR 1.137(a) or (b)) must be filed and granted to restore the application to pending status.

SEND ALL CORRESPONDENCE TO:

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SIGNATURE

James A. LaBarre

NAME

28,632

REGISTRATION NUMBER

Patent
Attorney's Docket No. 032326-132

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

In re Patent Application of)
Christophe CLAVIER et al) Group Art Unit: Unassigned
Application No.: 09/807,607) Examiner: Unassigned
Filed: April 16, 2001)
For: COUNTERMEASURE METHOD IN)
AN ELECTRONIC COMPONENT)
USING A SECRET KEY)
CRYPTOGRAPHIC ALGORITHM)

**SUBMISSION OF TRANSLATION TO ANNEX AND
SECOND PRELIMINARY AMENDMENT**

Assistant Commissioner for Patents
Washington, D.C. 20231

Sir:

Submitted herewith is a translation of the Annex to the International Preliminary Examination Report into English, comprising Claims 1-12. A check in the amount of \$130, to cover the surcharge under 37 C.F.R. §1.492(f), is also being submitted herewith.

Prior to examination and the calculation of filing fees, kindly amend the above-identified application as follows, to conform with the claims appearing in the Annex:

IN THE CLAIMS:

Kindly replace claims 1 8, 10 and 13, as follows.

1. A countermeasure method against attacks by differential analysis of current consumption in an electronic component using a cryptographic algorithm having a secret key, comprising the following steps:

executing some instructions in the algorithm that are critical to said attacks with a first manipulating means to deliver output data on the basis of input data, and executing others of said critical instructions with other manipulating means that are derived from said first manipulating means by complementation of at least one of said input data and said output data, so that the output data and data derived from said output data are unpredictable.

8. A countermeasure method according to claim 7, wherein said second manipulating means are such that, for the same input data, the complement of the output data of the first manipulating means is produced as output data, and wherein said second manipulating means are used in the second sequence for the fourteenth round.

10. A countermeasure method according to claim 1 wherein said manipulating means are tables of constants.

13. An electronic component which provides countermeasures against attacks on a secret key cryptographic algorithm, comprising:

a program memory having stored therein a plurality of different manipulating means for producing output data in response to input data;

a processor which executes instructions in said algorithm that are critical to said attacks, in accordance with a selected one of said manipulating means; and

means for generating a random value for selecting the manipulating means to be employed during a given execution of said algorithm, such that output data produced thereby is unpredictable.

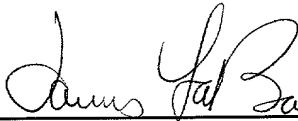
09307607 060404
FOI090 20920860

REMARKS

Entry of the foregoing amendment is respectfully requested. This amendment is intended to conform the claims with those appearing in the Annex to the International Preliminary Examination Report.

Respectfully submitted,

BURNS, DOANE, SWECKER & MATHIS, L.L.P.

By: 
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Date: June 1, 2001

032326-132-0001

Patent
Attorney's Docket No. 032326-132

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

In re Patent Application of)
)
Christophe CLAVIER et al) Group Art Unit: Unassigned
)
Application No.: Unassigned) Examiner: Unassigned
)
Filed: April 16, 2001)
)
For: COUNTERMEASURE METHOD IN)
AN ELECTRONIC COMPONENT)
USING A SECRET KEY)
CRYPTOGRAPHIC ALGORITHM)

PRELIMINARY AMENDMENT

Assistant Commissioner for Patents
Washington, D.C. 20231

Sir:

Prior to examination and the calculation of filing fees, kindly amend the above-identified application as follows:

IN THE SPECIFICATION:

Page 1, immediately following the title appearing on lines 1 and 2, insert the following:

--This disclosure is based upon, and claims priority from French Application No. 98/12989, filed on October 16, 1998 and International Application No. PCT/FR99/02172, filed September 13, 1999, which was published on April 27, 2000 in a language other than English, the contents of which are incorporated herein by reference.

Background of the Invention--

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PCT 0990-032326-132

Page 8, between lines 17 and 18, insert the following heading:

--Summary of the Invention--.

Page 9, between lines 21 and 22, insert the following heading:

--Brief Description of the Drawings--.

Page 10, between lines 17 and 18, insert the following heading:

--Detailed Description--.

IN THE CLAIMS:

Cancel claims 11 and 12.

Kindly replace claims 1-10, as follows.

1. (Amended) A countermeasure method for an electronic component using a cryptographic algorithm having a secret key, comprising the following steps:
 - executing critical instructions in the algorithm with a first manipulating means to deliver output data on the basis of input data, and
 - executing said critical instructions with other manipulating means so that the output data and data derived from said output data are unpredictable.
2. (Amended) A countermeasure method according to claim 1, wherein said first and said other manipulating means are selected for use on the basis of one-half probability statistical relationship.

3. (Amended) A countermeasure method according to claim 2, wherein said algorithm comprises sixteen computation rounds, and wherein said method comprises executing a first sequence and a second sequence, each of which is made up of at least the first three rounds, such that the order in which the sequences are executed is a function of the one-half probability statistical relationship, with the first sequence using the first manipulating means in each round, and the second sequence using the other manipulating means in at least the first round.

4. (Amended) A countermeasure method according to claim 3, wherein each of the first and second sequences is made up of the first three rounds.

5. (Amended) A countermeasure method according to claim 3, wherein said other manipulating means consist of second means such that, for the same input data, the complement of the output data of the first manipulating means is produced as output data.

6. A countermeasure method according to claim 2, wherein said algorithm comprises sixteen computation rounds, and wherein said method comprises executing a first sequence and a second sequence, each of which is made up of at least the last three rounds, such that the order in which the sequences are executed is a function of the one-half probability statistical relationship, with the first sequence using the first manipulating means in each round, and the second sequence using the other manipulating means.

7. (Amended) A countermeasure method according to claim 6, wherein each of the first and second sequences is made up of the last three rounds, and wherein the other manipulating means used in the second sequence comprise second manipulating means and a third manipulating means.

8. (Amended) A countermeasure method according to claim 6, wherein said second manipulating means are such that, for the same input data, the complement of the output data of the first manipulating means is produced as output data, and wherein said second manipulating means are used in the second sequence for the fourteenth round.

9. (Amended) A countermeasure method according to claim 8, wherein said third manipulating means are such that, for the complement of the input data, the complement of the output data of the first manipulating means is produced as output data, and wherein said third manipulating means are used in the second sequence for the fifteenth round and the sixteenth round.

10. (Amended) A countermeasure method according to claim 1 wherein said manipulating means are constants tables.

Add the following claims.

--13. (New) An electronic component which provides countermeasures against attacks on a secret key cryptographic algorithm, comprising:

a program memory having stored therein a plurality of different manipulating means for producing output data in response to input data;

a processor which executes critical instructions in said algorithm in accordance with a selected one of said manipulating means; and

means for generating a random value for selecting the manipulating means to be employed during a given execution of said algorithm, such that output data produced thereby is unpredictable.

14. (New) The electronic component of claim 13 wherein said manipulating means comprise tables of constants.

15. (New) The electronic component of claim 13 wherein said different manipulating means respectively produce sets of output data that are complementary to one another.

16. (New) The electronic component of claim 13, wherein said component is a smart card.--


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REMARKS

Entry of the foregoing amendment is respectfully requested. This amendment is intended to place the claims in a more conventional format and eliminate the multiple dependency of the claims.

Respectfully submitted,

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Date: April 16, 2001

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Attachment to Preliminary Amendment dated April 16, 2001

Marked-up Claims 1-10

1. (Amended) A countermeasure method [in] for an electronic component using a cryptographic algorithm having a secret key [(K), use of], comprising the following steps:

executing critical instructions in the algorithm [comprising the use of] with a first manipulating means [(TC₀)] to deliver output data [(S)] on the basis of input data [(E), the output data and/or derived data being manipulated by critical instructions, said countermeasure method being characterized in that it makes provision to use], and

executing said critical instructions with other manipulating means [(TC₁)] so that the output data and [the derived] data derived from said output data are unpredictable.

2. (Amended) A countermeasure method according to claim 1, [characterized in that the use of the various means (TC₀, TC₁) is managed by] wherein said first and said other manipulating means are selected for use on the basis of one-half probability statistical relationship.

3. (Amended) A countermeasure method according to claim 2, [the use of the] wherein said algorithm [comprising] comprises sixteen computation rounds [(T1, ..., T16), said countermeasure method being characterized in that it], and wherein said method comprises executing a first sequence [(SEQA)] and a second sequence [(SEQB)], each of which is made up of at least the first three rounds [(T1, T2, T3)], such that the order in

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Attachment to Preliminary Amendment dated April 16, 2001

Marked-up Claims 1-10

which the sequences are executed [being] is a function of the one-half probability statistical relationship, with the first sequence [(SEQA)] using the first manipulating means [(TC₀)] in each round, and the second sequence [(SEQB)] using the other manipulating means [(TC₁)] in at least the first round [(T1)].

4. (Amended) A countermeasure method according to claim 3, [characterized in that] wherein each of the first and second sequences is made up of the first three rounds [(T1, T2, T3)].

5. (Amended) A countermeasure method according to claim 3, [or 4, characterized in that the other] wherein said other manipulating means consist of second means [(TC₁)] such that, for the same input data, [(E), they deliver as output] the complement [(\bar{S})] of the output data [(S)] of the first manipulating means [(TC₀)] is produced as output data.

6. A countermeasure method according to claim 2, [the use of the] wherein said algorithm [comprising] comprises sixteen computation rounds [(T1, ..., T16), said countermeasure method being characterized in that it], and wherein said method comprises executing a first sequence [(SEQA')] and a second sequence [(SEQB')], each of which is made up of at least the last three rounds [(T14, T15, T16)], such that the order in which

Attachment to Preliminary Amendment dated April 16, 2001

Marked-up Claims 1-10

the sequences are executed [being] is a function of the one-half probability statistical relationship, with the first sequence [(SEQA')] using the first manipulating means [(TC₀)] in each round, and the second sequence [(SEQB')] using the other manipulating means [(TC₁, TC₂)].

7. (Amended) A countermeasure method according to claim 6, [characterized in that] wherein each of the first and second sequences is made up of the last three rounds, and [in that] wherein the other manipulating means used in the second sequence comprise second manipulating means [(TC₁)] and a third manipulating means [(TC₂)].

8. (Amended) A countermeasure method according to claim 6 [or 7, characterized in that the second means (TC₁)], wherein said second manipulating means are such that, for the same input data_a [(E), they deliver as output] the complement [(\bar{S})] of the output data [(S)] of the first manipulating means [(TC₀)], and in that the second] is produced as output data, and wherein said second manipulating means are used in the second sequence [(SEQB')] for the fourteenth round [(T14)].

9. (Amended) A countermeasure method according to claim 8, [characterized in that the third means (TC₂)] wherein said third manipulating means are such that, for the complement of the input data_a [(E), they deliver as output] the complement [(\bar{S})] of the

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Attachment to Preliminary Amendment dated April 16, 2001

Marked-up Claims 1-10

output data [(S)] of the first manipulating means [(TC₀) and they] is produced as output data, and wherein said third manipulating means are used in the second sequence for the fifteenth round and the sixteenth round [(T15, T16)].

10. (Amended) A countermeasure method according to [any preceding claim, characterized in that the various] claim 1 wherein said manipulating means are constants tables.

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**Attachment to Submission of Translation Annex and
Second Preliminary Amendment dated June 1, 2001**

Marked-up Claims 1, 8, 10 and 13

1. (Twice Amended) A countermeasure method [for] against attacks by
differential analysis of current consumption in an electronic component using a
cryptographic algorithm having a secret key, comprising the following steps:

executing [critical] some instructions in the algorithm that are critical to said
attacks with a first manipulating means to deliver output data on the basis of input data, and
executing others of said critical instructions with other manipulating means
that are derived from said first manipulating means by complementation of at least one of
said input data and said output data, so that the output data and data derived from said
output data are unpredictable.

8. (Twice Amended) A countermeasure method according to claim [6] 7,
wherein said second manipulating means are such that, for the same input data, the
complement of the output data of the first manipulating means is produced as output data,
and wherein said second manipulating means are used in the second sequence for the
fourteenth round.

10. (Twice Amended) A countermeasure method according to claim 1 wherein
said manipulating means are tables of constants [tables].

**Attachment to Submission of Translation Annex and
Second Preliminary Amendment dated June 1, 2001**

Marked-up Claims 1, 8, 10 and 13

13. (Amended) An electronic component which provides countermeasures against attacks on a secret key cryptographic algorithm, comprising:
- a program memory having stored therein a plurality of different manipulating means for producing output data in response to input data;
 - a processor which executes [critical] instructions in said algorithm that are critical to said attacks, in accordance with a selected one of said manipulating means; and
 - means for generating a random value for selecting the manipulating means to be employed during a given execution of said algorithm, such that output data produced thereby is unpredictable.

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NEW CLAIMS

1. Countermeasure method against attacks by differential analysis of current consumption in an electronic component, implementing a cryptographic algorithm having a secret key (K), the implementation of the algorithm comprising the use of first digital processing means (TC₀) to provide an output datum (S) from an input datum (E), the said output datum and/or data derived from this output datum being manipulated by instructions of the said algorithm which are critical in relation to the said attacks, **wherein** the countermeasure method provides for the use of other digital processing means (TC₁) in an alternative manner with the said first means, the said other means being obtained from the said first means by complementation of the input datum and/or of the output datum, such that the output datum and the said derived data are unpredictable.

2. Countermeasure method according to claim 1, **wherein** the use of the different means (TC₀, TC₁) is managed by a statistical distribution of probability one-half.

3. Countermeasure method according to claim 2, the implementation of the algorithm comprising sixteen runs of calculation (T1,...,T16), **wherein** it comprises the execution of a first sequence (SEQA) and of a second sequence (SEQB) formed by at least the first three runs (T1, T2, T3), the order of execution of the sequences being a function of the statistical distribution of probability one-half, the first sequence (SEQA) using the first means (TC₀) in each run, the second sequence (SEQB) using the other means (TC₁) in at least the first run (T1).

4. Countermeasure method according to claim 3, **wherein** the first and the second sequences are each formed from the first three runs (T1, T2, T3).

5. Countermeasure method according to claim 3 or 4, **wherein** the other means consist of second means (TC₁) such that for a same input datum (E), they provide as an output the complement (/S) of the output datum (S) of the first means (TC₀).

6. Countermeasure method according to claim 2, the implementation of the algorithm comprising sixteen runs of computation (T1, ..., T16), **wherein** it comprises the execution of a first sequence (SEQA') and of a second sequence (SEQB'), each formed by at least the three last runs (T14, T15, T16), the order of execution of the sequences being a function of the statistical distribution of probability one-half, the first sequence (SEQA') using the first means (TC₀) in each run, the second sequence (SEQB') using the other means (TC₁, TC₂).

7. Countermeasure method according to claim 6, **wherein** the first and the second sequences are each formed by the three last runs, and in that the other means used in the second sequence comprise second means (TC₁) and third means (TC₂).

8. Countermeasure method according to claim 6 or 7, **wherein** the second means (TC₁) are such that, for a same input datum (E), they provide at the output the complement (/S) of the output datum (S) of the first means (TC₀), and **wherein** the second means are used in the second sequence (SEQB') for the fourteenth run (T14).

9. Countermeasure method according to claim 8, **wherein** the third means (TC₂) are such that, for the complement of the input datum (E), they provide at the output the complement (/S) of the output datum (S) of the first means (TC₀) and are used in the second sequence, for the fifteenth run and the sixteenth run (T15, T16).

10. Countermeasure method according to any one of the preceding claims, **wherein** the various means are tables of constants.

11. Electronic component comprising a microprocessor, a program memory, and a working memory, permitting the implementation of a cryptographic algorithm having a secret key (K), first digital processing means (TC₀) being provided to provide an output datum (S) from an input datum (E), the said output datum and/or data derived from this datum being manipulated by instructions of the said algorithm which are critical in relation to the said

attacks by differential analysis of current consumption, **wherein** it comprises means for implementation of a countermeasure method against the said attacks according to any one of the preceding claims 1-10, comprising other, fixed means (TC_1) of digital processing with the first means in program memory of the said component, and means of generation of a random value (RND1) of 0 or 1 for managing the use of the said first means and other means (TC_0 , TC_1).

12. Smart card comprising an electronic component according to claim 11.

8/PKTS

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JC03 Rec'd PCT/PAJ

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A COUNTERMEASURE METHOD IN AN ELECTRONIC COMPONENT
USING SECRET-KEY CRYPTOGRAPHIC ALGORITHM

The present invention relates to a countermeasure method in an electronic component using a secret-key cryptographic algorithm. Such a component is used in applications in which access to services or to data is strictly controlled. It has an architecture formed around a microprocessor and around memories, including a program memory that contains the secret key.

Such components are used in particular in chip cards or "smart cards" for certain applications thereof. For example, such applications are to accessing certain databanks, to remote payment, e.g. for television, to dispensing gasoline, or to going through highway turnpikes or tolls.

Such components or such cards therefore use a secret-key cryptographic algorithm, the best known of which is the Data Encryption Standard (DES) algorithm. Other secret-key algorithms exist, such as the RC5 algorithm or the COMP128 algorithm. Naturally, this list is not exhaustive.

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5 In general and in short, the function of such algorithms is to compute an enciphered message on the basis of a message applied as input (to the card) by a host system (server, cash dispenser or automatic teller, etc.) and on the basis of the secret key contained in the card, and to send back to the host system the enciphered message, thereby making it possible, for example, for the host system to authenticate the component or the card, to interchange data, etc.

10 Unfortunately, it has become apparent that such components or such cards are vulnerable to attacks that consist in differentially analyzing current consumption and that enable ill-intentioned third parties to find the secret key. Such attacks are referred to as "Differential Power Analysis" ("DPA") attacks.

15 The principle of such attacks is based on the fact that the power consumption of (i.e. the current consumed by) the microprocessor executing the instructions varies depending on the bit that is manipulated.

20 In particular, an instruction of the microprocessor that manipulates a data bit generates two different current profiles depending on whether the bit is a "1" bit or a "0" bit. Typically, if the instruction manipulates a "0", then, at that execution instant, a first amplitude is obtained for the consumed current, and if the instruction manipulates a "1", a second amplitude is obtained for the consumed current, which amplitude is different from the first amplitude.

25 30 The characteristics of the cryptographic algorithms are well known: computations performed, and

parameters used. The only unknown is the secret key contained in the program memory. The secret key cannot be deduced merely from knowledge of the message applied as input and of the enciphered message sent back.

5 However, in a cryptographic algorithm, certain items of computed data depend only on the message applied in non-enciphered form to the input of the card, and on the secret key contained in the card. Other items of data computed in the algorithm can also
10 be recomputed only on the basis of the enciphered message (generally delivered in unenciphered form as output from the card to the host system) and of the secret key contained in the card. More precisely, each bit of these particular items of data can be determined
15 on the basis of the input or output message, and of a limited number of particular bits of the key.

 Thus, each bit of a particular item of data corresponds to a subkey formed by a particular group of bits of the key.

20 The bits of such particular items of data that can be predicted are referred to below as "target bits".

 The basic idea behind DPA attack is to use the difference in the current consumption profile of an instruction depending on whether it is manipulating a
25 "1" or a "0", and to use the possibility of a target bit being computed by the instructions of the algorithm on the basis of a known input or output message and of a hypothesis on the corresponding subkey.

 The principle of DPA attack is thus to test a
30 given subkey hypothesis by applying a selection Boolean function to a large number of current measurement curves, each relating to an input message known to the

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attacker, the selection Boolean function being a function of the subkey hypothesis and being defined for each curve by the value predicted for a target bit.

By making a hypothesis on the subkey in question, it is possible to predict the "0" or "1" value that the target bit will take for a given input or output message.

It is then possible for the predicted "0" or "1" value taken by the target bit for the subkey hypothesis in question to be applied as a selection Boolean function for sorting the curves into two packets. A first packet groups together the curves that have seen manipulation of the target bit at "0" and a second packet groups together the curves that have seen manipulation of the target bit at "1" depending on the subkey hypothesis. By averaging the current consumption in each packet, it is possible to obtain a mean consumption curve $M_0(t)$ for the first packet and a mean consumption curve $M_1(t)$ for the second packet.

If the subkey hypothesis is correct, the first packet really does group together all those curves from among the N curves which have seen the manipulation of the target bit at "0" and the second packet really does group together all those curves from among the N curves which have seen the manipulation of the target bit at "1". The mean consumption curve $M_0(t)$ of the first packet then has a mean consumption everywhere except at the moments at which the critical instructions are executed, with a current consumption characteristic of the manipulation of the bit at "0" ($profile_0$). In other words, for all of the curves, all of the manipulated bits have had the same chances of being equal to "0" as

of being equal to "1" except for the target bit which has always taken the value "0". This may be written as follows:

$$M_0(t) = [(profile_0 + profile_1)/2]_{t \neq t_{ci}} + [profile_0]_{t_{ci}}$$

i.e.

$$M_0(t) = [V_{m_t}]_{t \neq t_{ci}} + [profile_0]_{t_{ci}}$$

where t_{ci} represents the critical instants at which a critical instruction was executed.

Similarly, the mean consumption curve $M_1(t)$ of the second packet corresponds to a mean consumption everywhere except at the moments at which the critical instructions are executed, with a current consumption curve characteristic of the manipulation of the target bit at "1" ($profile_1$). The following can be written:

$$M_1(t) = [(profile_0 + profile_1)/2]_{t \neq t_{ci}} + [profile_1]_{t_{ci}}$$

i.e.

$$M_1(t) = [V_{m_t}]_{t \neq t_{ci}} + [profile_1]_{t_{ci}}$$

As described, the two profiles $profile_0$ and $profile_1$ are not equal. The difference in the curves $M_0(t)$ and $M_1(t)$ then gives a signal $DPA(t)$ whose amplitude is equal to $profile_0 - profile_1$ at the critical instants t_{ci} at which the critical instructions manipulating said bit are executed, i.e., in the example shown in Figure 1, at the places tc_0 to tc_6 , and whose amplitude is approximately equal to zero outside the critical instants.

If the subkey hypothesis is incorrect, the sorting does not correspond to reality. Statistically, there are then, in each packet, as many curves having really

seen the manipulation of the target bit at "0" as there are curves having seen the manipulation of the target bit at "1". The resulting mean curve $M0(t)$ is then situated around a mean value given by $(profile_0 + profile_1)/2 = V_m$, because, for each of the curves, all of the manipulated bits, including the target bit have as many chances of taking the value "0" as they have of taking the value "1".

The same reasoning on the second packet leads to a mean current consumption curve $M1(t)$ whose amplitude is situated around a mean value given by $(profile_0 + profile_1)/2 = V_m$

The signal $DPA(t)$ supplied by the difference $M0(t) - M1(t)$ is, in this case, substantially equal to zero. The signal $DPA(t)$ in the case of a incorrect subkey hypothesis is shown in Figure 2.

Thus, the DPA attack makes use of the difference in the current consumption profile while an instruction is being executed, depending on the value of the manipulated bit, to sort current consumption curves on the basis of a Boolean selection function for a given subkey hypothesis. By differentially analyzing the mean power consumption between the two resulting packets of curves, it is possible to obtain an information signal $DPA(t)$.

The procedure of a DPA attack thus consists substantially in the following:

a- drawing N random messages (e.g. N equals 1,000);

b- having the algorithm executed by the card for each of the N random messages, while recording the

current consumption curve each time (as measured on the power supply terminal of the component);

c- making a hypothesis on a subkey;

d- predicting, for each of the random messages;
5 the value taken by one of the target bits whose value depends only on the bits of the (input or output) message and on the subkey taken as a hypothesis, in order to obtain the Boolean selection function;

e- sorting the curves on the basis of this Boolean selection function (i.e. on the basis of the "0" or "1" value predicted for the target bit for each curve under the subkey hypothesis);

f- computing, in each packet, the resulting mean current consumption curve; and

g- determining the difference in said mean curves,
15 so as to obtain the signal $DPA(t)$.

If the hypothesis on the subkey is correct, the Boolean selection function is correct, and the curves of the first packet really do correspond to the curves for which the message applied as input or delivered as output has given a target bit at "0" in the card, and the curves of the second packet really do correspond to the curves for which the message applied as input or delivered as output has given a target bit of "1" in
20 the card.

This is the case shown in Figure 1: the signal $DPA(t)$ is therefore not zero at the instants tc_0 to tc_6 corresponding to the execution of the critical instructions (i.e. the instructions that manipulate the target bit).
25

It should be noted that the attacker does not need to know the critical instants with precision. It
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suffices for the attacker to have at least one critical instant in the acquisition period.

If the subkey hypothesis is incorrect, the sorting does not correspond to the reality and, in each packet, there are as many curves really corresponding to a target bit at "0" as there are curves corresponding to a target bit at "1". The $DPA(t)$ signal is substantially zero everywhere (case shown in Figure 2). It is necessary to return to step c- and a new subkey hypothesis must be made.

If the hypothesis proves to be correct, it is possible to move on to assessing other subkeys, until the key has been reconstructed to the maximum possible extent. For example, with a DES algorithm, a 64-bit key is used, only 56 bits of which are working bits. With a DPA attack, it is possible to reconstruct at least 48 bits of the 56 working bits.

An object of the present invention is to use a countermeasure method in an electronic component, which method gives a signal $DPA(t)$ that is zero even when the subkey hypothesis is correct.

In this way, nothing makes it possible to distinguish the case when the subkey hypothesis is correct from the case when the subkey hypothesis is incorrect. By means of this countermeasure, the electronic component is protected against DPA attacks.

In the invention, the countermeasure method makes it possible to make the target bits, i.e. the data manipulated by critical instructions, unpredictable.

Because of the countermeasure, for each message applied as input, a target bit manipulated by a critical instruction has an equal probability of taking

the value 0 or the value 1. In each packet of curves that the attacker makes under a given subkey hypothesis, by means of the Boolean selection function that the attacker has computed, there are as many
5 curves that have really manipulated "0" target bits as curves that have really manipulated "1" bits. The signal $DPA(t)$ is always zero, whether the subkey hypothesis is correct or incorrect.

As characterized, the invention therefore provides
10 a countermeasure method in an electronic component using a cryptographic algorithm having a secret key, embodiment of the algorithm comprising the use of first means to deliver output data on the basis of input data, the output data and/or derived data being
15 manipulated by critical instructions. According to the invention, the countermeasure method makes provision to use other means so that the output data and the derived data are unpredictable.

According to the invention, the use of the various
20 means is managed by a one-half probability statistical relationship.

Other characteristics and advantages of the invention are described in more detail in the following description given by way of non-limiting example and
25 with reference to the accompanying drawings, in which:

Figures 1 and 2 (described above) show the $DPA(t)$ signal that can be obtained as a function of a hypothesis on a subkey of the secret key K , in a DPA attack;

30 Figures 3 and 4 are flow charts showing how the first rounds and the last rounds of the DES algorithm are executed;

Figure 5 is a block diagram of the SBOX operation used in the DES algorithm;

Figure 6 shows an example of a one-input and one-output elementary table of constants used in the SBOX operation;

Figures 7 and 8 show examples of flow charts of execution of the first and last rounds of the DES algorithm, in an embodiment of the countermeasure method of the invention;

Figures 9 and 10 respectively show second and third elementary tables of constants of the invention;

Figure 11 is a general flow chart of the execution of the DES in an embodiment of the countermeasure method of the invention; and

Figure 12 is a simplified block diagram of a smart card including an electronic component in which the countermeasure method of the invention is used.

The present invention is explained below in an example of application to the DES cryptographic algorithm. The invention is not limited to this example alone. It is applicable to secret-key cryptographic algorithms in general.

The DES cryptographic algorithm (referred to below as the "DES" or as the "DES algorithm") comprises 16 computation rounds, referenced T1 to T16, as shown in Figures 3 and 4.

The DES begins with an initial permutation IP on the input message M (Figure 3). The input message M is a 64-bit word f. After permutation, a 64-bit word e is obtained that is cut in two to form the input parameters L0 and R0 of the first round (T1). L0 is a word d having 32 bits and containing the 32 most

significant bits of the word e. R0 is a word h having 32 bits and containing the 32 least significant bits of the word e.

5 The secret key K, which is a 64-bit word g, itself undergoes permutation and compression to deliver a word r having 56 bits.

10 The first round comprises an operation EXP PERM on the parameter R0, which operation consists in expansion and permutation so as to deliver as output a word l having 48 bits.

15 This word l is combined with a parameter K1 in an EXCLUSIVE OR type operation referenced XOR so as to deliver a word b having 48 bits. The parameter K1 which is a 48-bit word m is obtained from the word r by shifting by one position (operation referenced SHIFT in Figures 3 and 4) followed by permutation and by compression (operation referenced COMP PERM).

20 The word b is applied to an operation referenced SBOX at the output of which a word a having 32 bits is obtained. This particular operation is explained in more detail with reference to Figures 5 and 6.

The word a undergoes a permutation P PERM, giving as output the 32-bit word c.

25 This word c is combined with the input parameter L0 of the first round T1, in a logic EXCLUSIVE OR type operation, referenced XOR, which delivers as output the 32-bit word g.

30 The word h (=R0) from the first round supplies the input parameter L1 of the following round (T2), and the word g from the first round supplies the input parameter R1 of the following round. The word p from

the first round supplies the input \underline{r} of the following round.

The other rounds T2 to T16 proceed similarly except as regards the SHIFT operation which is performed on one or two positions depending on the round in question.

Each round T1 thus receives as input the parameters L_{i-1} , R_{i-1} , and \underline{r} and delivers as output the parameters L_i , R_i , and \underline{r} for the following round T_{i+1} .

At the end of the DES algorithm (Figure 4), the enciphered message is computed on the basis of the parameters L_{16} and R_{16} delivered by the last round T16.

The enciphered message C comprises, in practice, the following operations:

- forming a 64-bit word e' by inverting the positions of the words L_{16} and R_{16} , and then by concatenating them; and

- applying the permutation IP^{-1} that is the inverse of the permutation at the beginning of the DES, so as to obtain the 64-bit word f' forming the enciphered message C.

The SBOX operation is shown in detail in Figures 5 and 6. It comprises a constants table TC_0 for delivering output data \underline{a} as a function of input data \underline{b} .

In practice, this constants table TC_0 is in the form of eight elementary constants tables TC_{01} to TC_{08} , each receiving as input only 6 bits of the word \underline{b} , so as to deliver as output only 4 bits of the word \underline{a} .

Thus, the elementary constants table TC_{01} shown in Figure 6 receives as input data the bits b_1 to b_6 of the word \underline{b} and delivers as output the bits a_1 to a_4 of the word \underline{a} .

In practice, these eight elementary constants tables are stored in a program memory of the electronic component.

5 In the SBOX operation of the first round T1, a particular bit of the output data a of the constants table TC₀ depends on only 6 bits of the data b applied as input, i.e. on only 6 bits of the secret K and of the input message (M).

10 In the SBOX operation of the last round T16, a particular bit of the output data a of the constants table TC₀ can be recomputed on the basis of only 6 bits of the secret key K and of the enciphered message (C).

15 Going back to the principle of the DPA attack, if the target bit is chosen to be a bit of the output data a, it suffices to make a hypothesis on 6 bits of the key K in order to predict the value of a target bit for a given input message (M) or for a given output message (C). In other words, for the DES, it suffices to make a hypothesis on a 6-bit subkey.

20 In a DPA attack on such an algorithm for a given target bit, it is necessary to distinguish a correct subkey hypothesis from 64 possible hypothesis.

25 Thus, by taking only eight bits from the word a as target bits (one output bit per elementary constants table TC₀1 to TC₀8), it is possible to discover up to $6 \times 8 = 48$ bits of the secret key by making DPA attacks on each of the target bits.

30 In the DES, instructions that are critical for DPA attacks are to be found at the beginning of the algorithm and at the end of the algorithm.

At the beginning of the DES algorithm, the items of data that can be predicted on the basis of an input

message M and of a subkey hypothesis are the items of data a and g computed in the first round (T1).

The data a of the first round T1 (Figure 3) is the output data of the SBOX operation of the round in question. The data g is computed on the basis of the data a, by permutation (P PERM) and EXCLUSIVE OR operation with the input parameter L0.

The data c of the first round is data derived from the data a of the first round. The derived data c corresponds to a single permutation of bits of the data a.

The data l of the second round is data derived from the data g of the first round, because it corresponds to a permutation of the bits of the word g, certain bits of the word g further being duplicated.

Knowing a and g, it is possible also to know said derived data.

The critical instructions of the beginning of the algorithm are the critical instructions that manipulate either the data that can be predicted or the data a of the first round, or else derived data.

The critical instructions manipulating the data a of the first round T1 or the derived data c are thus instructions of the end of the SBOX operation, of the P PERM operation, and of the beginning of the XOR operation of the first round T1.

The critical instructions manipulating the data g or derived data are all the instructions from the end of the XOR operation of the first round T1 to the instructions of the start of the SBOX operation of the second round T2, and the instructions of the start of

the XOR operation of the end of the end of the third round T3 ($L2 = h(T2) = g(T1)$).

At the end of the DES algorithm, the items of data that can be predicted on the basis of an enciphered message C and of a subkey hypothesis are the data a of the sixteenth round T16 and the data L15 equal to the word h of the fourteenth round T14.

The critical instructions manipulating the data a of the sixteenth round or derived data are the instructions of the sixteenth round for the end of the SBOX operation, for the P PERM permutation operation and for the start of the XOR operation.

For the data L15, the critical instructions manipulating said data or the derived data are all the instructions from the end-of-XOR-operation instructions of the fourteenth round T14 to the start-of-SBOX-operation instructions of the fifteenth round T15, and the start-of-XOR-operation instructions of the end the sixteenth round T16.

The countermeasure method of the invention as applied to this DES algorithm consists in having, for each critical instruction, as many chances that the critical instruction is manipulating an item of data as that it is manipulating the complement of said item of data. Thus, regardless of the target bit on which the DPA attack can be made, there are as many chances that the critical instructions that manipulate the bit are manipulating a "1" as that they are manipulating a "0".

In practice, this must apply for each of the potential target bits: in other words, since the attacker has the choice between a plurality of possible attacks, i.e. between a plurality of possible Boolean

selection functions for performing the curve sorting, for any given subkey hypothesis, the embodiment of the countermeasure method of the invention must make sure that the data manipulated by each of the critical instructions takes randomly, and every other time, a value or the complement of that value. As regards applying the countermeasure method of the invention to the DES algorithm, it is thus necessary to apply the countermeasure to the critical instructions at the beginning of the DES, and to the critical instructions at the end of the DES, in order to be fully protected.

In the DES, all of the data manipulated by critical instructions is output data or data derived from output data of an SBOX operation.

At the beginning of the DES, the data that can be predicted is the data a and g of the first round T1. The data a is the output data of the SBOX operation of the first round. The data g is computed on the basis of the data a, since $g = P \text{ PERM}(a) \text{ XOR } L0$. g is thus data derived from the output data a of the SBOX operation of the first round. Thus, all of the data manipulated by the critical instructions of the beginning of the DES result directly or indirectly from the output data a of the SBOX operation of the first round.

As regards the end of the DES, the data that can be predicted is the data a of the sixteenth round T16 and the data g of the fourteenth round T14, g being equal to L15.

The data a is the output data of the SBOX operation of the sixteenth round T16.

As regards the data L15, it is computed, during normal execution of the DES algorithm, on the basis of the output data a of the SBOX operation of the fourteenth round T14: $L15 = P \text{ PERM } (a) \text{ XOR } L14$.

5 If the items of output data a of the particular SBOX operations are made unpredictable, all of the items of derived data are also made unpredictable: thus all of the items of data manipulated by the critical instructions of the DES algorithm are made
10 unpredictable.

The SBOX operation thus corresponds to first means which consist of a constants table TC_0 , and which are used in each round to deliver output data E on the basis of input data S.

15 An embodiment of the countermeasure method as applied to the DES algorithm may consist in using at least one other constants table as other means for making the output data a unpredictable, so that all of the items of said output data and/or of derived data
20 manipulated by the critical instructions are unpredictable.

In executing the algorithm, the use of the various means, i.e., in the example, of the various constants tables is managed on the basis of a one-half
25 probability statistical relationship.

The other constants table or the other constants tables are such that they make the complemented item of data correspond to one and/or the other of the items of input data d and of output data of the first constants
30 table TC_0 .

Figures 7 and 8 thus show an embodiment of the countermeasure of the invention as applied to the DES algorithm.

Figure 7 shows the beginning of the algorithm.
 5 The data operations that are not modified by the countermeasure method of the invention bear the same references as in Figure 3 described above.

At the beginning of the DES algorithm, a second constants table TC_1 is provided in the SBOX operation of the first round T_1 . All of the items of data affected by the second constants table TC_1 are given a ' sign or a $\bar{}$ sign in these figures. It can be seen that the critical instructions of the beginning of the DES manipulate all of the items of data affected by the countermeasure method.
 10
 15

It can be observed that, since the first constants table is made up of eight first constants tables, the second constants table is also made up of eight second constants tables.

20 In the embodiment shown, the first constants table TC_0 and the second constants table TC_1 are such that, for the same input data E , the second table delivers as output the complement \bar{S} of the output data S delivered by the first table.

25 Figure 9 shows such a second elementary table TC_{11} delivering output that is complemented relative to the first elementary table TC_{01} shown in Figure 6.

With such a second constants table TC_1 , the complement \bar{a} of the data a obtained with the first constants table TC_0 is obtained at the output the SBOX operation of the first round T_1 . Similarly, in the
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first round T1, the complemented data \bar{g} is obtained, and, in the second round T2, the complemented items of data \bar{h} , $\bar{L2}$, \bar{l} , and \bar{b} are obtained.

By using the first table or the second table to deliver the output data on the basis of a one-half probability statistical relationship, all of the potential target bits at the start of the DES that are manipulated by the critical instructions have as many chances of taking the value "1" as they do of taking the value "0".

At the end of the DES algorithm, using the countermeasure method of the invention requires the use of a plurality of constants tables that are different from the first table because it is necessary to consider both the output data \underline{a} computed in the fourteenth round T14, and the output data \underline{a} computed in the sixteenth round T16 in order to make all of the items of data manipulated by the critical instructions of the end of the DES unpredictable.

An embodiment of the countermeasure method as applied to the end of the DES algorithm is shown in Figure 8.

It makes provision to use two constants tables TC_1 and TC_2 .

In the SBOX operation of the fourteenth round T14, the second constants table TC_1 (already used for the beginning of the DES) is used.

And a third constants table TC_2 is used in the SBOX operations of the fifteenth and sixteenth rounds.

This third constants table TC_2 is such that it supplies the complement \bar{S} of the output data S to the

complement \bar{E} of the input data E of the first constants table TC_0 . An example of a corresponding third elementary constants table TC_2 , on the basis of the first elementary constants table TC_0 , is shown in Figure 10.

By using such constants tables, it appears in Figure 8 that all of the critical instructions manipulate complemented data.

The invention is not limited to these examples of constants tables TC_1 and TC_2 . Other possibilities exist. For example, for the countermeasure method applied at the end of the DES, it is also possible to combine use of the constants table TC_1 with use of another constants table defined relative to the first constants table TC_0 as supplying the output data S to the complement \bar{E} of the input data.

In general, the end of the DES requires the use of different constants tables, as a function of the rounds in question, so that all of the items of data manipulated by the critical instructions of the end of the DES are unpredictable.

The embodiment described with reference to Figures 7 and 8 does have a drawback however: the countermeasure method applied at the input of the DES produces intermediate computed results $L3'$ and $R3'$ which are incorrect. All of the following intermediate results are therefore incorrect as well.

Similarly, at the end of the DES, the countermeasure method applied at the input of the DES produces intermediate computed results $L16'$ and $R16'$ that are incorrect.

In all cases, the enciphered message is incorrect.

In this embodiment of the invention, it is thus necessary to make provision to be capable of resuming the remainder of the algorithm with the correct intermediate results every time, once the critical instructions have been passed.

In practice, since the critical instructions of the start of the DES are to be found in the first three rounds, these three rounds are duplicated. In other words, provision is made to execute two sequences, each of which comprises at least the first three rounds T1, T2, T3. A first sequence SEQA uses the first constants table TC₀ in each round. The other sequence SEQB uses the second constants table TC₁ at least in the first round T1. In the example shown, the first constants table is used in the two following rounds T2 and T3.

As described, in the countermeasure method of the invention, the use of the various means, i.e. in the example, the use of the various constants tables, is managed on the basis of a one-half probability statistical relationship. This one-half probability statistical relationship is then more particularly applied to the order in which the various means are used, i.e., in the example, to the order in which the two sequences SEQA and SEQB are executed.

Similarly, to have the correct parameters L16 and R16 at the end of the DES for forming the enciphered message C, the three rounds T14, T15, and T16 (Figure 7) that contain the end-of-DES critical instructions are also duplicated. Two sequences are thus executed, each of which comprises at least the last three rounds T14, T15, T16. A first sequence SEQA' uses the first

constants table TC_0 in each round. The other sequence SEQB' uses the other constants tables TC_1 and TC_2 . As above, the one-half probability statistical relationship is then applied to the order in which the two sequences SEQA' and SEQB' are executed.

The critical instructions are then executed twice, once in each sequence. But, at the time at which any of the critical instructions of either one of the sequences is executed, the probability of manipulating an item of data is equal to the probability of manipulating its complement.

The computation program of the DES that is used in the electronic component must therefore be modified to include the countermeasure method of the invention. An example of an execution flow chart of the invention, using the countermeasure method at the start of and at the end of the DES, as in the embodiment described with reference to Figures 7 and 8 is shown in Figure 11. In this example, each of the sequences SEQA and SEQB comprise the first three rounds and each of the sequences SEQA' and SEQB' comprise the last three rounds.

The computation program then consists mainly, at the beginning of computation, in saving the input parameters referenced DATAIN and KEY which, in practice, correspond to the parameters L_0 , R_0 , and \underline{r} , in a temporary memory zone referenced CONTEXT0.

In this computation program, a first loop counter FR is then positioned at 0, and a value RND1 equal to 0 or to 1 is drawn randomly.

In the example, if RND1 is equal to 1, the sequence SEQB of T_1 , T_2 , T_3 is performed first, in

which sequence the second constants table TC_1 is used in round T1, and the first table TC_0 is used for rounds T2 and T3. The output parameters $L3'$, $R3'$ (which have incorrect values) are saved in a temporary memory zone referenced CONTEXT2.

If FR is not equal to 1, it is put to 1, the input parameters are retrieved from CONTEXT0, and the value of RND1 is complemented. In the example, RND1=0 is obtained. The other sequence SEQA of T1, T2, T3 is then executed in which the first constants table is used in all three rounds T1, T2, and T3. The output parameters (correct values) are saved in a temporary memory zone referenced CONTEXT1.

If FR is at 1, both sequences have been performed. CONTEXT1 is then retrieved to deliver the intermediate results $L3$, $R3$ having the correct values, in the next round (T4).

If RND1 is equal to zero, the procedure begins with T1(TC_0), T2(TC_0), T3(TC_0) and finishes with T1(TC_1), T2(TC_0), T3(TC_0).

At the end of the round T13, the parameters $L13$ and $R13$ delivered by this round are saved in the temporary memory CONTEXT0, and the procedure continues for the remaining rounds T14, T15, and T16 in a manner similar to the first rounds.

In all cases, the number of instructions must be exactly the same regardless of the computation path. This is why, in particular, in the application example described, provisions is made also to save the incorrect values ($L3'$, $R3'$, or $L16'$, $R16'$) in the temporary memory zone CONTEXT2.

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If any difference exists between the two possible paths, there is then a possibility of fruitful DPA attack.

The countermeasure method of the invention is not limited to the particular embodiment described with reference to the algorithm DES. It is applicable to any secret-key cryptographic algorithm. In general, for any embodiment of an algorithm comprising the use of first means for delivering output data on the basis of input data, with the output data and/or derived data being manipulated by critical instructions, the countermeasure method of the invention comprises the use of other means, so that the output data and the derived data are unpredictable.

The use of the various means, i.e. of the first means and of the other means, is managed using a one-half probability statistical relationship.

The other means may comprise a plurality of different means. They are such that they make the complemented item of data correspond to one or other of the items of input data or of output data of the first means.

In the example described wherein the countermeasure method is applied to DES, the first means consist of the first constants table TC_0 . The other means consist, at the start of the DES, of the second constants table TC_1 . At the end of the DES, they consist of two different constants tables TC_1 and TC_2 , in the example.

In order to apply the countermeasure method of the invention to a given secret-key cryptographic algorithm, it is necessary firstly to determine all

those items of data of this algorithm which can be predicted, and all those critical instructions (i.e. instructions that are critical in the DPA attack sense) which manipulate such items of data or derived items of data. It is then necessary to identify in the algorithm first means and other means in the sense of the invention, so that all of the data manipulated by the critical instructions is unpredictable. For the DES algorithm, the first means are the constants table TC₀. The other means are, for example, other constants tables. These means may be different operations for other algorithms. For the same algorithm, the means may consist of different operations depending on the identified critical instructions.

The electronic component 1 using such a countermeasure method in a secret-key cryptographic algorithm typically comprises, as shown in Figure 12, a microprocessor μ P, a program memory 2, and a working memory 3. To be capable of managing the use of the various means of the invention, which are, in the example described, the various constants tables stored in the program memory, means 4 for generating a random 0 or 1 value are provided that, with reference to Figure 11, deliver the value of RND1 each time the DES is executed. Such a component may particularly be used in a smart card CP to improve its tamper-resistance.

CLAIMS

1/ A countermeasure method in an electronic component using a cryptographic algorithm having a secret key (K), use of the algorithm comprising the use of first means (TC₀) to deliver output data (S) on the basis of input data (E), the output data and/or derived data being manipulated by critical instructions, said countermeasure method being characterized in that it makes provision to use other means (TC₁) so that the output data and the derived data are unpredictable.

2/ A countermeasure method according to claim 1, characterized in that the use of the various means (TC₀, TC₁) is managed by a one-half probability statistical relationship.

3/ A countermeasure method according to claim 2, the use of the algorithm comprising sixteen computation rounds (T₁, ..., T₁₆), said countermeasure method being characterized in that it comprises executing a first sequence (SEQA) and a second sequence (SEQB), each of which is made up of at least the first three rounds (T₁, T₂, T₃), the order in which the sequences are executed being a function of the one-half probability statistical relationship, the first sequence (SEQA) using the first means (TC₀) in each round, and the second sequence (SEQB) using the other means (TC₁) in at least the first round (T₁).

4/ A countermeasure method according to claim 3, characterized in that each of the first and second sequences is made up of the first three rounds (T₁, T₂, T₃).

5/ A countermeasure method according to claim 3 or 4, characterized in that the other means consist of second means (TC_1) such that, for the same input data (E), they deliver as output the complement (\bar{S}) of the output data (S) of the first means (TC_0).

6/ A countermeasure method according to claim 2, the use of the algorithm comprising sixteen computation rounds (T_1, \dots, T_{16}), said countermeasure method being characterized in that it comprises executing a first sequence (SEQA') and a second sequence (SEQB'), each of which is made up of at least the last three rounds (T_{14}, T_{15}, T_{16}), the order in which the sequences are executed being a function of the one-half probability statistical relationship, the first sequence (SEQA') using the first means (TC_0) in each round, the second sequence (SEQB') using the other means (TC_1, TC_2).

7/ A countermeasure method according to claim 6, characterized in that each of the first and second sequences is made up of the last three rounds, and in that the other means used in the second sequence comprise second means (TC_1) and third means (TC_2).

8/ A countermeasure method according to claim 6 or 7, characterized in that the second means (TC_1) are such that, for the same input data (E), they deliver as output the complement (\bar{S}) of the output data (S) of the first means (TC_0), and in that the second means are used in the second sequence (SEQB') for the fourteenth round (T_{14}).

9/ A countermeasure method according to claim 8, characterized in that the third means (TC_2) are such that, for the complement of the input data (E), they

deliver as output the complement (\bar{S}) of the output data (S) of the first means (TC₀) and they are used in the second sequence for the fifteenth round and the sixteenth round (T15, T16).

5 10/ A countermeasure method according to any preceding claim, characterized in that the various means are constants tables.

10 11/ An electronic component using the countermeasure method according to any preceding claim, said electronic component being characterized in that the various means (TC₀, TC₁, TC₂) for delivering output data on the basis of input data are fixed in a program memory of said component, and in that said component includes means for generating a random 0 or 1 value
15 (RND1) for managing the use of said various means.

 12/ A smart card including an electronic component according to claim 11.

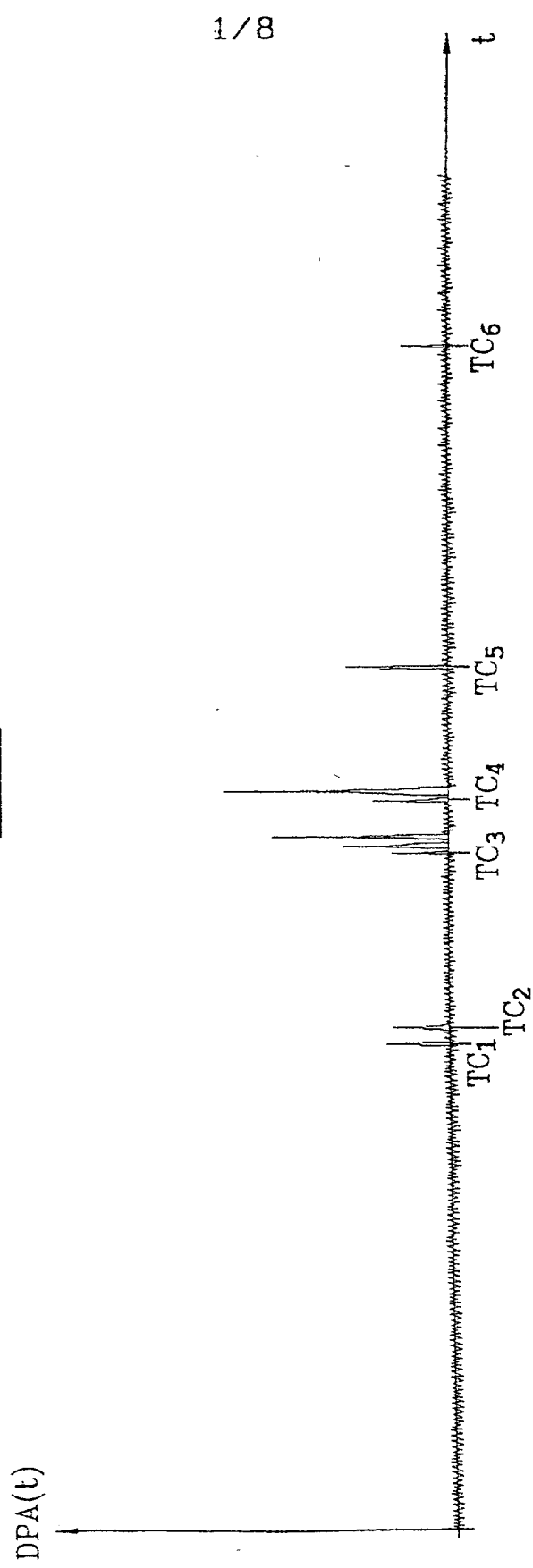
0900709090

ABSTRACT

In an electronic component using a cryptographic algorithm having a secret key (K), use of the algorithm comprising the use of first means (TC₀) to deliver output data (S) on the basis of input data (E), the output data and/or derived data being manipulated by critical instructions, a countermeasure method makes provision to use other means (TC₁ and/or TC₂) so that the output data and the derived data are unpredictable.

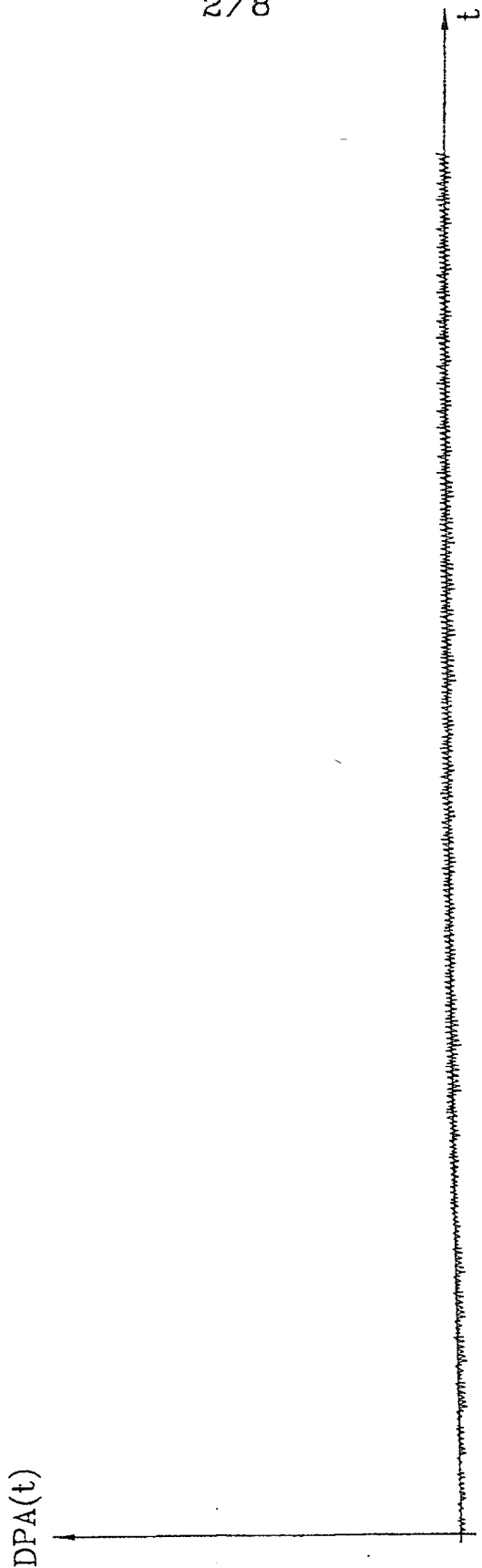
09807607 080101

FIG.1

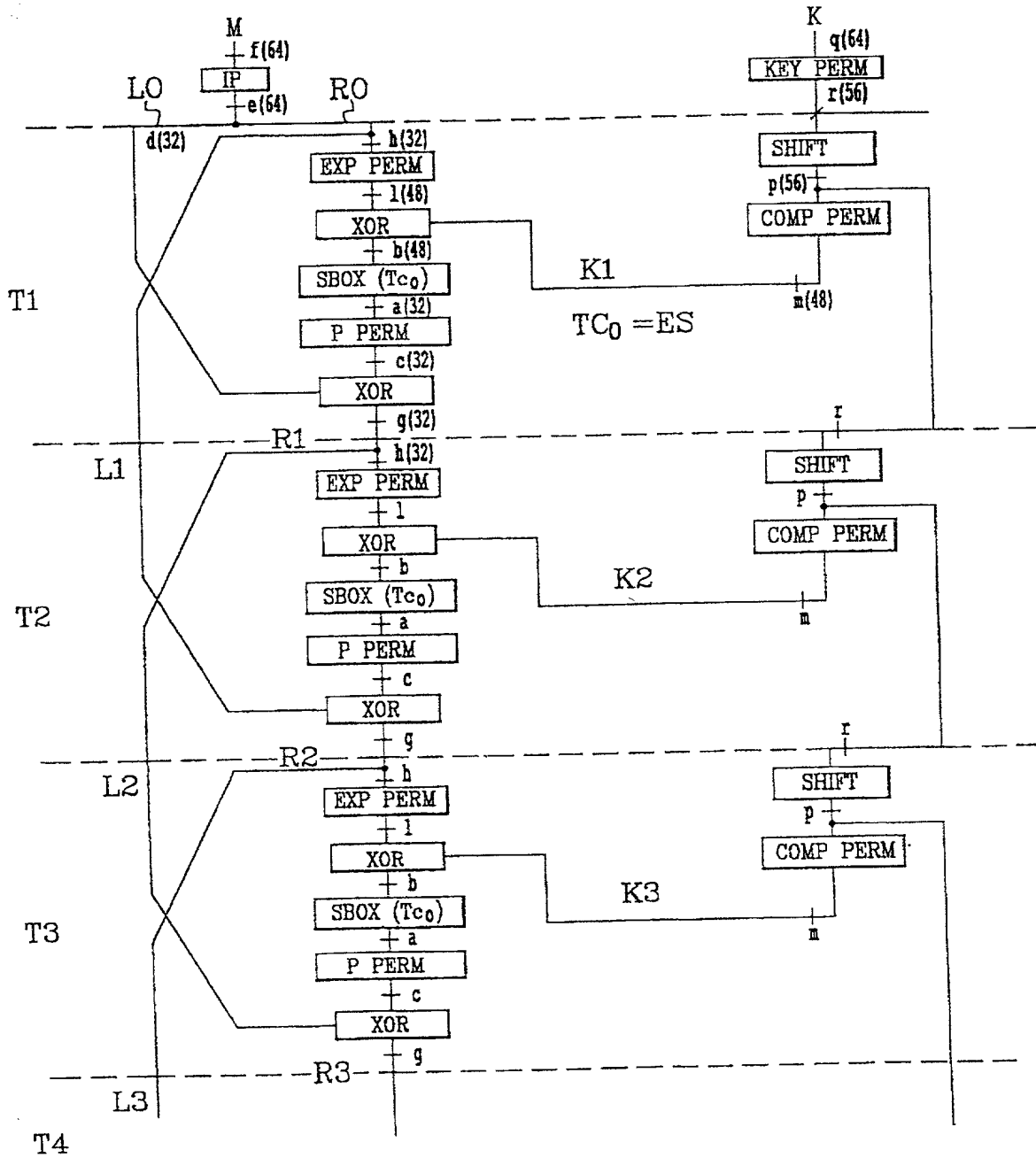


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FIG.2



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FIG.3

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T13

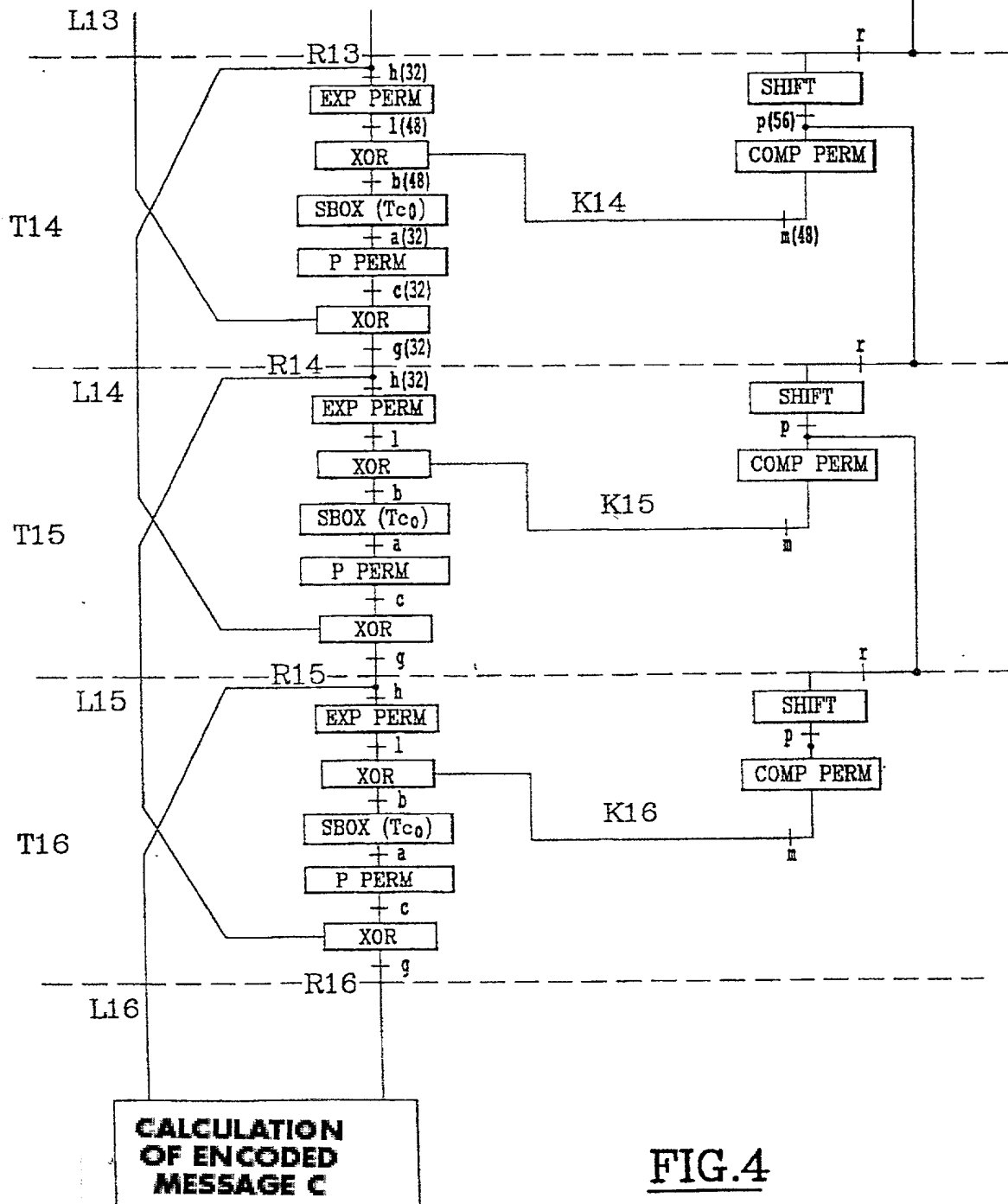
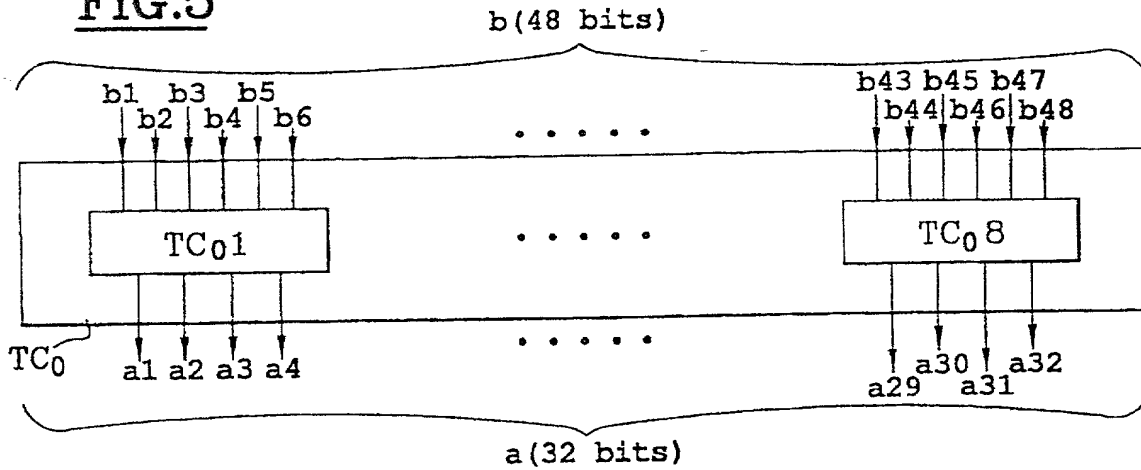


FIG.4

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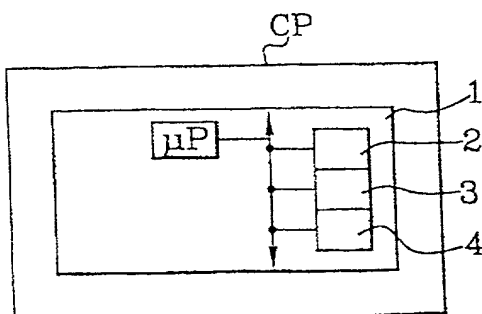
FIG.5**FIG.6**

TC ₀₁	E1=b1b2b3b4b5b6	S1=a1a2a3a4
	000000	1101
	000001	0101
	⋮	⋮
	111111	1010

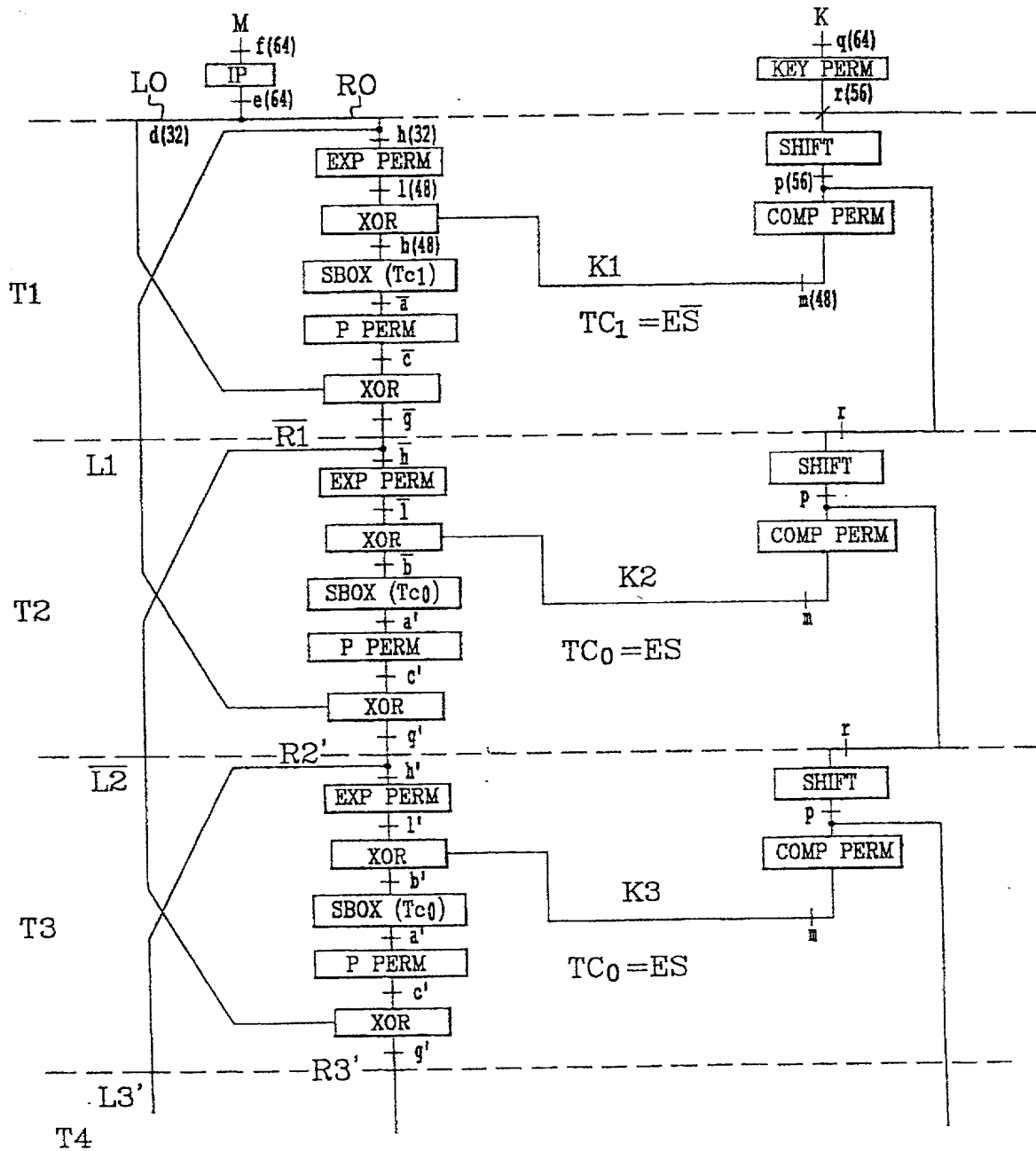
E1=b1b2b3b4b5b6	/S1=a1a2a3a4	TC ₁₁
000000	0010	
000001	1010	
⋮	⋮	
111111	0101	

FIG.9

TC ₂₁	/E1=b1b2b3b4b5b6	/S1=a1a2a3a4
	000000	0101
	⋮	⋮
	111110	1010
	111111	0010

FIG.10**FIG.12**

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FIG.7

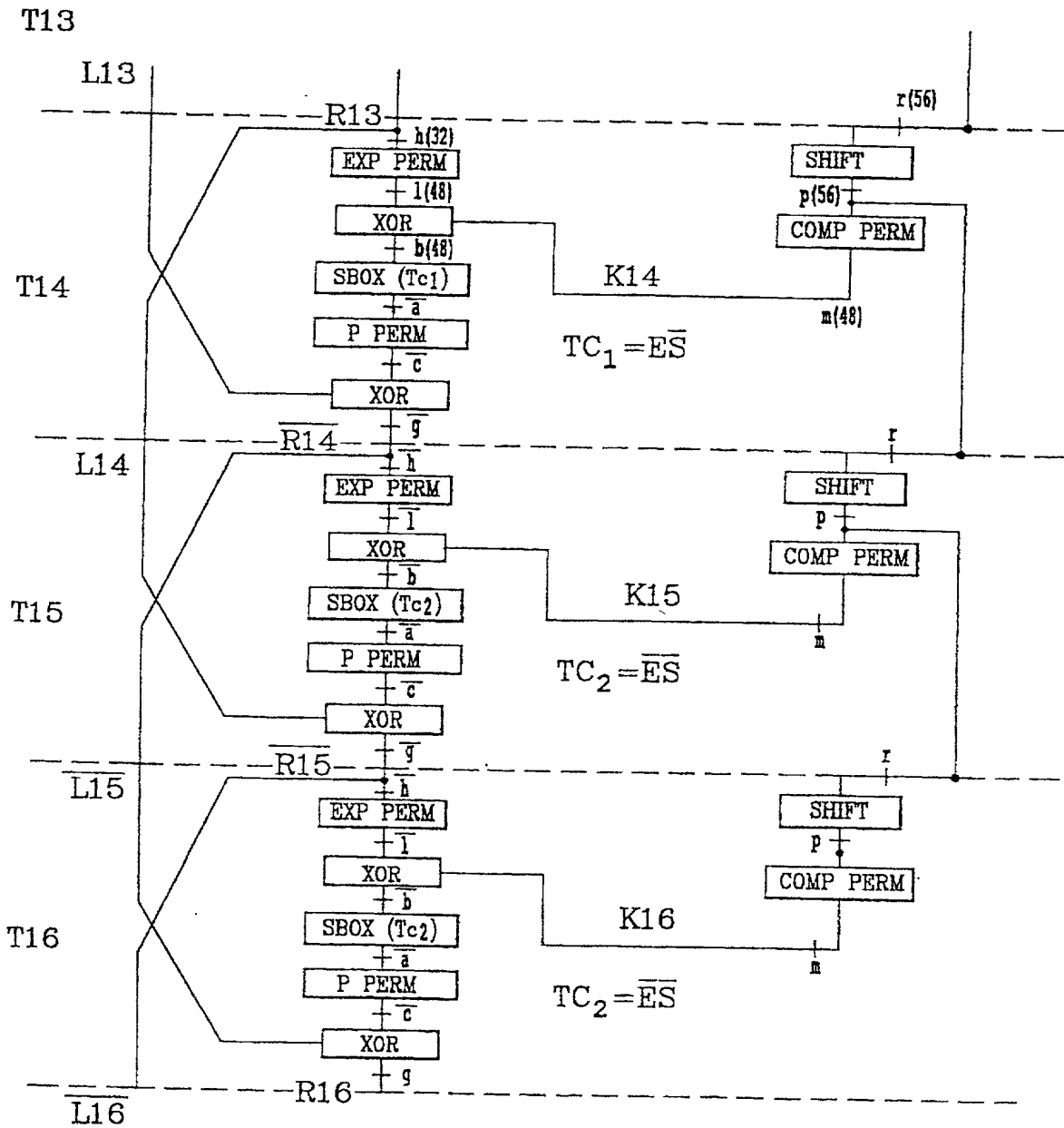


FIG.8

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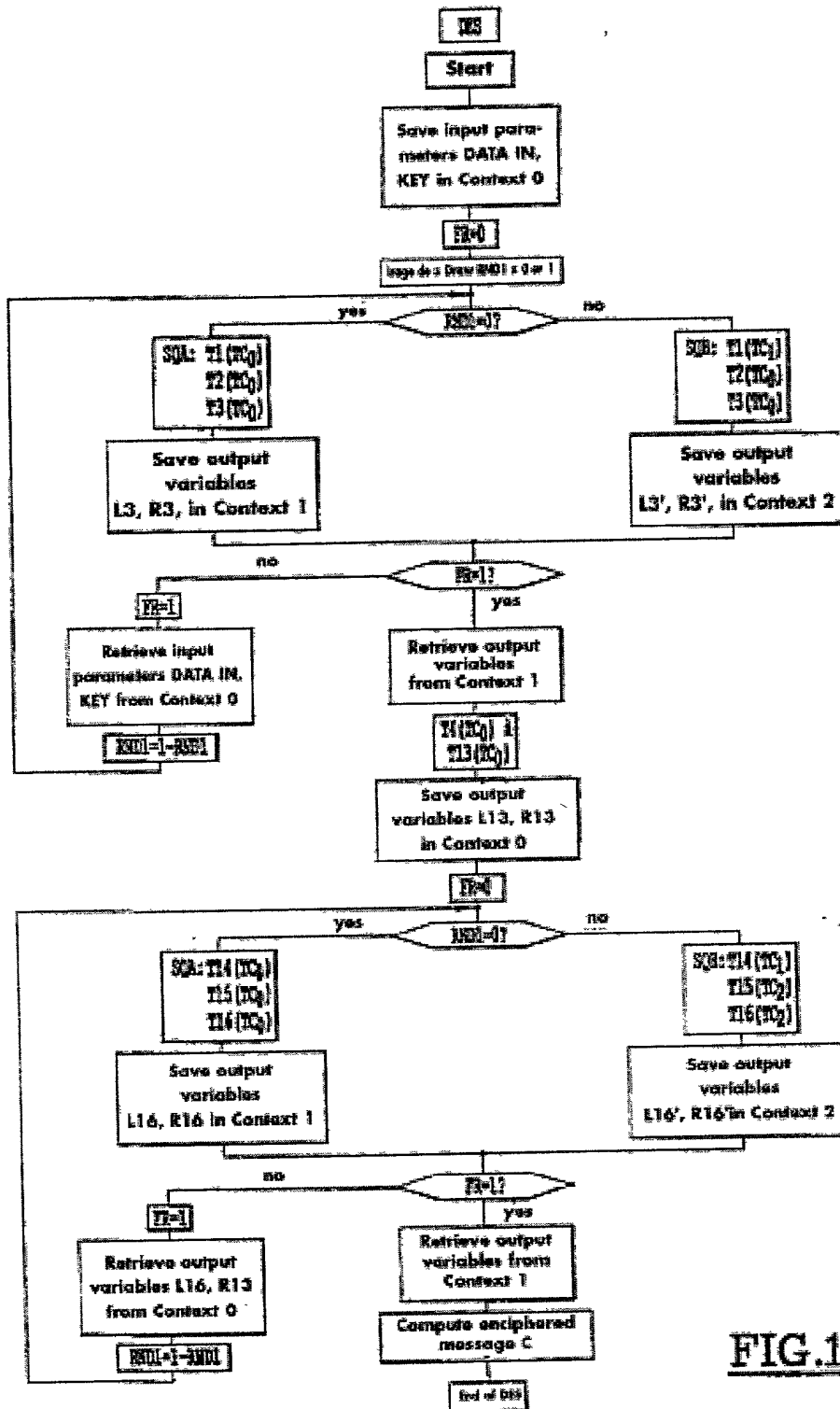


FIG. 11

COMBINED DECLARATION FOR PATENT APPLICATION AND POWER OF ATTORNEY
(Includes Reference to Provisional and International (PCT) Applications)

Attorney's Docket No.
032326-132

As a below named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated below next to my name;

I BELIEVE I AM THE ORIGINAL, FIRST AND SOLE INVENTOR (IF ONLY ONE NAME IS LISTED BELOW) OR AN ORIGINAL, FIRST AND JOINT INVENTOR (IF PLURAL NAMES ARE LISTED BELOW) OF THE SUBJECT MATTER WHICH IS CLAIMED AND FOR WHICH A PATENT IS SOUGHT ON THE INVENTION ENTITLED:

Countermeasure Method in an Electronic Component Using a Secret Key Cryptographic Algorithm

The specification of which (check only one item below):

- ☐ is attached hereto.
- ☐ was filed as United States Patent Application Number 09/807,607
on April 16, 2001
and was amended on _____ (if applicable).
- ☐ was filed as International (PCT) Application Number _____
on _____
and was amended on _____ (if applicable).

I HAVE REVIEWED AND UNDERSTAND THE CONTENTS OF THE ABOVE-IDENTIFIED SPECIFICATION, INCLUDING THE CLAIMS, AS AMENDED BY ANY AMENDMENT REFERRED TO ABOVE.

I ACKNOWLEDGE THE DUTY TO DISCLOSE TO THE U.S. PATENT AND TRADEMARK OFFICE ALL INFORMATION KNOWN TO ME TO BE MATERIAL TO PATENTABILITY AS DEFINED IN TITLE 37, CODE OF FEDERAL REGULATIONS, Sec. 1.56 (as amended effective March 16, 1992);

I do not know and do not believe the said invention was ever known or used in the United States of America before my or our invention thereof, or patented or described in any printed publication in any country before my or our invention thereof or more than one year prior to said application; that said invention was not in public use or on sale in the United States of America more than one year prior to said application; that said invention has not been patented or made the subject of an inventor's certificate issued before the date of said application in any country foreign to the United States of America on any application filed by me or my legal representatives or assigns more than six months prior to said application;

I hereby claim foreign priority benefits under Title 35, United States Code, §§ 119 (a)-(e) of any foreign application(s) for patent or inventor's certificate or of any International (PCT) Application(s) designating at least one country other than the United States of America listed below and have also identified below any foreign application(s) for patent or inventor's certificate or any PCT International (PCT) Application(s) designating at least one country other than the United States of America filed by me on the same subject matter having a filing date before that of the application(s) of which priority is claimed:

PRIOR FOREIGN/PCT APPLICATION(S) AND ANY PRIORITY CLAIMS UNDER 35 U.S.C. §119:

COUNTRY (If PCT, indicate "PCT")	APPLICATION NUMBER	DATE OF FILING (day, month, year)	PRIORITY CLAIMED UNDER 35 U.S.C. §119
France	98/12989	16 October 1998	<input type="checkbox"/> Yes <input type="checkbox"/> No

I hereby claim the benefit under Title 35, United States Code § 119(e) of any United States provisional application(s) listed below.

(APPLICATION NUMBER)	(FILING DATE)
(APPLICATION NUMBER)	(FILING DATE)

COMBINED DECLARATION FOR PATENT APPLICATION AND POWER OF ATTORNEY (CONT'D)
(Includes Reference to Provisional and International (PCT) Applications)

Attorney's Docket
No. 032326-132

I hereby claim the benefit under Title 35, United States Code, § 120 of any United States application(s) or International (PCT) Application(s) designating the United States of America that is/are listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in that/those prior application(s) in the manner provided by the first paragraph of Title 35, United States Code, § 112, I acknowledge the duty to disclose to the U.S. Patent and Trademark Office all information known to me to be material to the patentability as defined in Title 37, Code of Federal Regulations § 1.56, which became available between the filing date of the prior application(s) and the national or international filing date of this application:

PRIOR U.S. APPLICATIONS OR INTERNATIONAL (PCT) APPLICATIONS DESIGNATING THE U.S. FOR BENEFIT UNDER 35 U.S.C. § 120:

U.S. APPLICATIONS		STATUS (check one)		
U.S. APPLICATION NUMBER	U.S. FILING DATE	PATENTED	PENDING	ABANDONED
		<input type="checkbox"/>	<input type="checkbox"/>	<input type="checkbox"/>
PCT APPLICATIONS DESIGNATING THE U.S.				
PCT APPLICATION NO.	PCT FILING DATE	U.S. APPLICATION NUMBERS ASSIGNED (if any)		
FR99/02172	13 September 1999			

I hereby appoint the following attorneys and agent(s) to prosecute said application and to transact all business in the U.S. Patent and Trademark Office connected therewith and to file, prosecute and to transact all business in connection with international applications directed to said invention:

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21839

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21839

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I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

COMBINED DECLARATION FOR PATENT APPLICATION AND POWER OF ATTORNEY (CONT'D) (Includes Reference to Provisional and International (PCT) Applications)	Attorney's Docket No. 032326-132
---	-------------------------------------

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RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS)		
FULL NAME OF FIFTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS)		
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FULL NAME OF SEVENTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
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POST OFFICE ADDRESS (HOME ADDRESS)		
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RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS)		
FULL NAME OF NINTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
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POST OFFICE ADDRESS (HOME ADDRESS)		
FULL NAME OF TENTH JOINT INVENTOR, IF ANY	SIGNATURE	DATE
RESIDENCE (CITY & STATE/COUNTRY)		CITIZENSHIP
POST OFFICE ADDRESS (HOME ADDRESS)		